

Introduction to **Distributed Ledger Technologies**

(the classy name of blockchains)

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iCU3E

What are we? WIN 222 NO 222 NO 222 39-0 0 **0** 3%-0 0 0







































no single point of failure





no single point of failureno central Authority





- no single point of failureno central Authority
- no need to trust the others





- no single point of failureno central Authority
- no need to trust the others

I want to add some data

















iCU3E Where is the data stored ? How do I read data ? How do I write data ? I want to add some data How do I make sure my data is not tempered ? How do I make sure what I read is correct? How do I become rich?





Talk Chain





Bitcoin PoW Consensus

"Old-style" Consensus Algorithms

Bitcoin PoW Consensus










Each process starts with an input value



Each process starts with an input value



Agreement

Old-style" Consensus Algorithms Bitcoin PoW Consensus Process (can be faulty) Channel (can be slow) Definition of Consensus:

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Agreement

Every correct process must agree on the same value.

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"Old-style" Consensus Algorithms Bitcoin PoW Consensus Bitcoin PoW Consensus Bitcoin PoW Consensus

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Definition of Consensus:

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Every correct process must agree on the same value.

Validity

If all processes propose the same value v, then all correct processes decide v.

Bitcoin PoW Consensus



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Termination

Bitcoin PoW Consensus



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iCU3E

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Bitcoin PoW Consensus



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iCU3E

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Integrity

Bitcoin PoW Consensus



Definition of Consensus:

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Every correct process must agree on the same value.

Validity

ICU3E

If all processes propose the same value v, then all correct processes decide v.

Termination

Every correct process decides some value.

Integrity

Every correct process decides at most one value, and if it decides some value *v*, then *v* must have been proposed by some process.





Impossibility results



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[FLP 1985] "Impossibility of distributed consensus with one faulty process"



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Costly algorithms



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Every node should be connected to every node, broadcast everything, and waits for n/2 responses



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Still puzzle researchers' minds



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Still puzzle researchers' minds [Bramas 2016] Packet efficient implementation of the omega failure detector



Impossibility results

[FLP 1985] <u>"Impossibility of distributed consensus with one faulty process"</u> [CAP theorem] <u>"Brewer's conjecture and the feasibility of consistent, available,</u> <u>partition-tolerant web services"</u>

Costly algorithms

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Costly algorithms

Every node should be connected to every node, broadcast everything, and waits for n/2 responses

Still puzzle researchers' minds [Bramas 2016] Packet efficient implementation of the omega failure detector

There exists consensus algorithms that tolerates up to n/3 Byzantins processes



So what's the problem with an algorithm tolerating n/3 byzantin nodes ?





So what's the problem with an algorithm tolerating n/3 byzantin nodes ?







"old-style" consensus algorithms work well when the participants are known, and not too many (eg, Hashgraph blockchain, 39 known participants in the "council")





We need some kind of



We need some kind of Proof



We need some kind of **Proof** of "existence"

We need some kind of Proof

of "existence"

iCU3E

that does not depend on how many nodes you control in the network

We need some kind of Proof

of "existence"

iCU3E

that does not depend on how many nodes you control in the network

but depends on ...



How much work can you do ?





How much work can you do ?

Your importance in the network depends on how many times you can execute a hash function. (no interest in simulating many nodes, because your hashing power will remain the same)





Basic principle of the Bitcoin Protocol :



Basic principle of the Bitcoin Protocol :

- Choose randomly one node






Nxt PoS Consensus

In more details

What is the database ? The blockchain

E



It's a data structure with one function: append(Block)

(you cannot remove a block, so to "remove" a data just overwrite it)



Who stores the data? Every nodes



Who stores the data? Every nodes

How to read data? just read all the blocks in order



Who stores the data? Every nodes

How to read data? just read all the blocks in order

Exemple 1: what's the value of v ?



E



Who stores the data? Every nodes

How to read data? just read all the blocks in order

Exemple 1: what's the value of v ?



Exemple 2: How much does Alice have ? and Bob ?





How to write data ?



Nxt PoS Consensus

In more details

How to write data ?

Send my data to everyone

E





Algorithms

How to write data ?

Send my data to everyone

The node that will be elected will include it in its block (hopefully)





Algorithms

How to write data ?

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The node that will be elected will include it in its block (hopefully)

Exemple : Oscar to Alice 10\$





3E

Algorithms

How to write data ?

- Send my data to everyone
- The node that will be elected will include it in its block (hopefully)

Exemple : Oscar to Alice 10\$



https://blockchain.info/unconfirmed-transactions







Each node includes "Oscar to Alice : 10\$" to their block



Each node includes "Oscar to Alice : 10\$" to their block

One Block is "randomly" selected to be appended to the blockchain



One Block is "randomly" selected to be appended to the blockchain







Distribution:Alice to Oscar : 50\$Alice : 100\$Alice to Daniel : 30\$













→ E

In more details

Exemple : Oscar to Alice 10\$



















Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

How to prevent anyone from appending a block? How to make sure data is not removed?

https://anders.com/blockchain/

4080782: I worked hard for this presentation

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

BE



Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

BE



It's easy for Alice to replace a block with another one:



Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

BE



It's easy for Alice to replace a block with another one:



But the other nodes will prefer the longest blockchain

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

3E



It's easy for Alice to replace a block with another one:



But the other nodes will prefer the longest blockchain

The deeper the block you modify, the harder it gets to generate a blockchain longer than the existing one



What if two nodes generates two different block at the same time ?





What if two nodes generates two different block at the same time ?



Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

SE

Algorithms

hash: #0000067 nonce: 3452 data: Bob to Alice: 10\$ prev: #032	
hash: #0000042 nonce: 98423 data: Bob to Alice: 1\$ prev: #032	

parts of the networks will receive this first and will try to append a block to this

(they don't want to change their choice when they receive the other one because that would mean that they wasted their computing power on this one)

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BE Bitcoin PoW Consensus Nxt PoS Consensus Algorithms E In more details hash: #0000091 hash: #0000067 nonce: 3452 same as before nonce: 3452 data: ... data: Bob to Alice: 10\$ prev: #0000067 prev: #032 hash: #0000012 hash: #0000042 same as before nonce: 3452 nonce: 98423 data: ... data: Bob to Alice: 1\$ prev: #0000042 prev: #032

SE Bitcoin PoW Consensus Algorithms Nxt PoS Consensus In more details hash: #0000091 hash: #0000067 nonce: 3452 If it finds the next nonce: 3452 block first data: ... data: Bob to Alice: 10\$ prev: #0000067 prev: #032 hash: #0000042 they receive the other block, stop working on nonce: 98423 extending #0000042 and start extending #0000091 (they don't want to continue with #0000042 because it becomes data: Bob to Alice: 1\$ harder to have a longer chain, so trying to extends here is a waste of their computing power)

prev: #032

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This block is orphaned


Nxt PoS Consensus



In more details

When finding a block, a node receives some coins

(a fixed reward + transaction fees chosen freely by the sender)



Nxt PoS Consensus



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Algorithms **BE**

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

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Oscar has 10\$ in this version of the blockchain

Algorithms 3

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

When finding a block, a node receives some coins

(a fixed reward + transaction fees chosen freely by the sender)



Oscar has 10\$ in this version of the blockchain

So the goal of the nodes, is to find a block first, and that this block is not orphaned.

This does not mean "extends only the version where I have the greatest balance" because you want to avoid wasting hashing power on a chain that has great chance to be orphaned

Nxt PoS Consensus

→ E

In more details

3E

Algorithms

What is the balance of Alice ?



Algorithms

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

Block Validation

BE







In more details

Double spending



Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign A -> B (1\$)



Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign A -> B (1\$)

but Alice also sign A -> A' (5\$) and send this one to the network



Nxt PoS Consensus



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Nxt PoS Consensus



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Double spending

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Now, no one will ever include A -> B (1\$) transaction because it's in conflict with A -> A' (5\$)



Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign A -> B (1\$)

Bob waits to see this transaction in a block before giving the sandwich



Nxt PoS Consensus



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Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign A \rightarrow B (1\$)

Bob waits to see this transaction in a block before giving the sandwich



In this chain, B never received the coins and A -> B (1\$) cannot be included anymore

Alice also sign A -> A' (5\$) and use a lots of power to extends this chain faster than the other



Nxt PoS Consensus

→ E

In more details

Double spending

Alice buy a sandwich to Bob and sign A -> B (1\$)

Bob waits to see this transaction in a block and waits for 10 confirmation blocks



Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign $A \rightarrow B$ (1\$)

Bob waits to see this transaction in a block and waits for 10 confirmation blocks





Nxt PoS Consensus



In more details

Double spending

Alice buy a sandwich to Bob and sign A -> B (1\$)

Bob waits to see this transaction in a block and waits for 10 confirmation blocks



generate 10 blocks on her own







8

Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

Double spending: If an adversary owns more than ½ the network hashing power, it can always perform a double spend



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Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

Double spending: If an adversary owns more than ½ the network hashing power, it can always perform a double spend

 \triangleright Orphaned blocks \Rightarrow wasted hashing power \Rightarrow reduce protocol security



Nxt PoS Consensus



In more details

- Double spending: If an adversary owns more than ½ the network hashing power, it can always perform a double spend
- \triangleright Orphaned blocks \Rightarrow wasted hashing power \Rightarrow reduce protocol security
 - The easier the PoW the more forks will occur. (the average time between two nodes to find a block should be much greater than the latency in the network)

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Nxt PoS Consensus



In more details

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Algorithms

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Bitcoin protocol :

- new blocks should arrive in average every 10 minutes
- reward : 12.5BTC (divided by two every 4 years)







Bitcoin protocol : new blocks should arrive in average every 10 minutes

H : hashrate, hash/s of the network







In more details

- H : hashrate, hash/s of the network
- p : probability of finding a block with 1 hash



Nxt PoS Consensus



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- $\frac{1}{p}$: expecting number of hash before finding a block



Nxt PoS Consensus



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- $\frac{1}{pH}$: expecting number of seconds before finding a block



Bitcoin PoW Consensus

Nxt PoS Consensus



In more details

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We want
$$\frac{1}{pH} = 600$$
 $p = \frac{1}{H600}$



Bitcoin PoW Consensus





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 2^{224} 1 1

$$p = \frac{2}{D} \times \frac{1}{2^{256}} = \frac{1}{D2^{32}}$$



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Bitcoin PoW Consensus



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 $D\,$ is the difficulty. The greater the difficulty, the smaller the probability $\,p\,$

$$D = \frac{H600}{2^{32}}$$









Fun facts

- Estimated Satoshi's balance : 1M BTC = 10 billion USD
- Current difficulty : 4,022,059,196,164 Current Hashrate : 29,911,310,546 GH/s (3 * 10^16 hash / s)

How many BTC I gifted to my cousin in 2013,


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https://digiconomist.net/bitcoin-energy-consumption



How to avoid wasting energy for leader election ?



How to avoid wasting energy for leader election ?







The more coins you own, the more chance you get of being elected





The more coins you own, the more chance you get of being elected

Election :





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Election :

For each address, take t = hash(adr + last block hash) / balance





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Election :

- For each address, take t = hash(adr + last block hash) / balance
- \triangleright The address with the smallest *t* is elected.





Using "old-style" Consensus algorithm





Using "old-style" Consensus algorithm

Elects n validators (with PoS for instance)

The validators use "old-style" Consensus algorithm to agree on the next blocks







Alice want to sell a singleplayer game



Alice want to sell a singleplayer game

But users should buy a licence to play it



Alice want to sell a singleplayer game

But users should buy a licence to play it

The contract is simple: a user pay 10\$ and he can play it



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IOTA BlockDAG

What is a smart contract ?

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SE

IOTA BlockDAG

What is a smart contract ?

Alice want to sell a singleplayer game

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"

"

var usersWhoOwnTheGame: Set

If an address *adr* sent to Alice 10\$ and add "buyTheGame()" in the data of a block **then** usersWhoOwnTheGame.insert(*adr*)



Alice add this to a block:

"

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var usersWhoOwnTheGame: Set

If an address *adr* sent to Alice 10\$ and add "buyTheGame()" in the data of a block **then** usersWhoOwnTheGame.insert(*adr*)

Then in addition to remember the balance of each address, each nodes remember the value of *usersWhoOwnTheGame* and updates it after each block if the condition is met.



If an address *adr* sent to Alice 10\$ and add "buyTheGame()" in the data of a block **then** usersWhoOwnTheGame.insert(*adr*)

Then in addition to remember the balance of each address, each nodes remember the value of *usersWhoOwnTheGame* and updates it after each block if the condition is met.

"



Alice add this to a block:

```
pragma solidity ^0.4.18;
contract AliceGame {
  mapping(address => bool) public usersWhoOwnTheGame;
  function buy() public payable {
    if(msg.value >= 10)
    {
      usersWhoOwnTheGame[msg.sender] = true;
    }
  }
}
```

test here: http://remix.ethereum.org/

-> gets translated to EVM instructions



<u>3e</u>

IOTA BlockDAG

PUSH1 0x80 PUSH1 0x40 **MSTORE CALLVALUE** DUP1 ISZERO PUSH2 0x10 JUMPI PUSH1 0x0 DUP1 REVERT JUMPDEST POP PUSH2 0x166 DUP1 PUSH2 0x20 PUSH1 0x0 CODECOPY PUSH1 0x0 RETURN STOP **PUSH1 0x80** PUSH1 0x40 **MSTORE** PUSH1 0x4 CALLDATASIZE





Each and **every** node in the network executes **all** the smart contracts, and keep track of all the values of all the variables.



Each and **every** node in the network executes **all** the smart contracts, and keep track of all the values of all the variables.

When a transaction arrives with data @*AliceGame.buy()* the node has to load the context of the smart contract into memory, executes the function, check if everything is ok, performs the corresponding actions.



Each and **every** node in the network executes **all** the smart contracts, and keep track of all the values of all the variables.

When a transaction arrives with data @*AliceGame.buy()* the node has to load the context of the smart contract into memory, executes the function, check if everything is ok, performs the corresponding actions.

A client can ask a node "what's the value of this variable @AliceGame.usersWhoOwnTheGame[playerAdr]"





Executing a contract has a cost depending on the complexity of the algorithm.


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In bitcoin, a transaction just consumes space in the block (fees encourage miners to include a transaction)



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When calling a contracts you declare ▶ The called function

How much gas at most the miner will spend executing it



Executing a contract has a cost depending on the complexity of the algorithm.

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In Ethereum, the cost of a calling a contract is calculated in Gas (each EVM instruction has a Gas Cost)

When calling a contracts you declare

- The called function
- How much gas at most the miner will spend executing it
- How much Ether you're willing to pay per Gas spent



Fun fact



Fun fact

You can buy this crypto cat for 150 ETH (10000\$)





Ethereum Smart Contracts

IOTA BlockDAG

Better Fact

https://www.dash.org/network/#section-governance





It does not scale



It does not scale

Bitcoin : 4 transactions / seconds



It does not scale

Bitcoin : 4 transactions / seconds Visa: 2 000 transactions / seconds





Sharding (each node own a fraction of the database)



Sharding (each node own a fraction of the database)

Side-chains



Sharding (each node own a fraction of the database)

Side-chains

DAG instead of Chain





Each transaction is a small block that reference two previous ones



Each transaction is a small block that reference two previous ones



IOTA BlockDAG

Each transaction is a small block that reference two previous ones



You come up with a DAG (Directed Acyclic Graph)

IOTA BlockDAG

Each transaction is a small block that reference two previous ones



You come up with a DAG (Directed Acyclic Graph)

You're only limited by bandwidth and storage



Each transaction is a small block that reference two previous ones





Each transaction is a small block that reference two previous ones





Each transaction is a small block that reference two previous ones



A new site and its parents should not create conflicts.



IOTA BlockDAG

The Tangle (IOTA)

How to read a value?



How to read a value?

If you take a tip, you can order transactions and do the same as in a blockchain





IOTA BlockDAG

The Tangle (IOTA)

How to read a value?

What if tips are conflicting?



A new site cannot confirm conflicting sites



Tip Selection Algorithm (TSA):

- so we know how to read values
- so we know where to extend the Tangle



IOTA BlockDAG

The Tangle (IOTA)



- so we know how to read values
- so we know where to extend the Tangle

In Bitcoin, we read values from, and we try to extend, the longest chain. If you don't follow this, you'll lose money.





In the Tangle, forks are ok if not conflicting

In the Tangle, forks are ok if not conflicting

But conflicting forks are worst in this case

In the Tangle, forks are ok if not conflicting

IOTA BlockDAG

But conflicting forks are worst in this case



In the Tangle, forks are ok if not conflicting So its better to have something like this






The Tangle (IOTA)

Should be chosen with higher probability





The Tangle (IOTA)

Compute cumulative weight to each site



nart Contracts

The Tangle (IOTA)

Compute cumulative weight to each site



IOTA BlockDAG Infinite Scalability nart Contracts The Tangle (IOTA) Compute cumulative weight to each site Perform a random walk

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nart Contracts IOTA BlockDAG

The Tangle (IOTA)

How many tips are left behind ?



IOTA BlockDAG

ICU3E

QuarkChain white paper :



https://quarkchain.io/QUARK%20CHAIN%20Public%20Version%200.3.4.pdf



QuarkChain white paper :

The QuarkChain Network is helping move blockchain into the next generation by increasing the current TPS capacity several-thousand fold of what it is now, to a projected about 100,000 TPS. The network being built is project to be free of congestion, making it affordable for all usage scenarios that demand speed and volume. We envision such a network applied to industries that demand higher TPS. Ultimately, the QuarkChain Network aims to build a high-throughput network to support applications such asdistributed social media, high frequency trading, Internet of Things (IoT), gaming, and payment.





Take Away











Lots of interesting open problems about :



Lots of interesting open problems about : Network impacts on blockchain security



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Network impacts on blockchain security

Properly defining the attacks



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Properly defining the attacks

▶ How to defend against them



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Also :

How to store efficiently blockchain data



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And also :



Lots of interesting open problems about :

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- How to defend against them

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And also :

What programs could run on the blockchain for distributed democracy / authority / management system for IoT devices



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Thank you for your attention !

Quentin Bramas. The Stability and the Security of the Tangle. 2018. $\langle hal-01716111v2 \rangle$